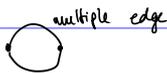


HW2

5. lots of you had trouble with the induction.



etc.

Note that the $k=2$ case is false! So multiple edges are not allowed. Loops are not allowed in \mathcal{W} either since they're cycles.

Closed walks of length 1 can be ignored since they're loops. So assume the walk has at least two steps.

Direct: $w = v_1 \dots v_k$, $v_1 = v_k$ (closed)

Find $i < j$ st $v_i = v_j$

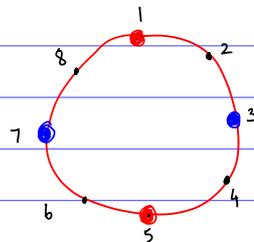
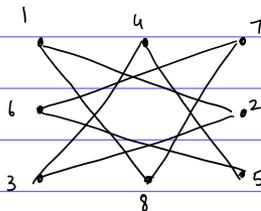
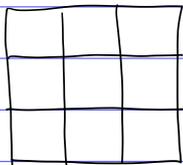
the # of vertices between v_i and v_j on w is

From your Etextbook defn 1.1.15: A cycle is a graph with $|V(\mathcal{G})| = |E(\mathcal{G})|$, and whose vertices can be placed around a circle, whose vertices are adjacent iff they appear consecutively on a circle.

No vertices are repeated between v_i and v_j because by definition this is the minimal closed walk. If there is more than one vertex between v_i and v_j , then this is a cycle of length at least 3.

So there must be exactly one vertex z between v_i and v_j . This edge must be repeated twice.

9. Most of you got this right.



This is a proof by picture. The graph of moves is an 8 cycle. It's easy to see that whenever two knights ^{of the same color} are in adjacent corners, the other knight must be sandwiched between them in one of the sites $\{2, 4, 6, 8\}$ which is not a corner.

